



To Adapt or Not to Adapt, That is the Ski Question

Building Online Adaptive Filters for Robust Performance

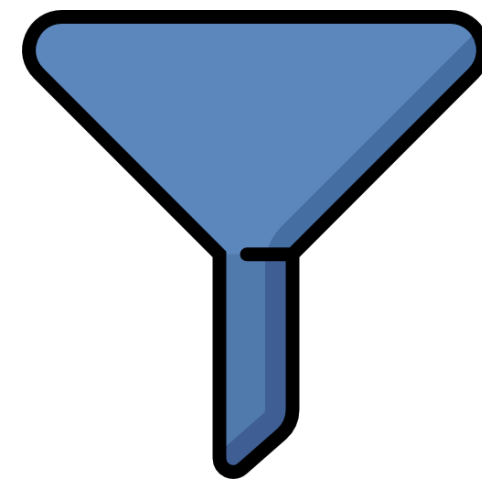
SIGMOD 2026

Yuvaraj Chesetti and Prashant Pandey



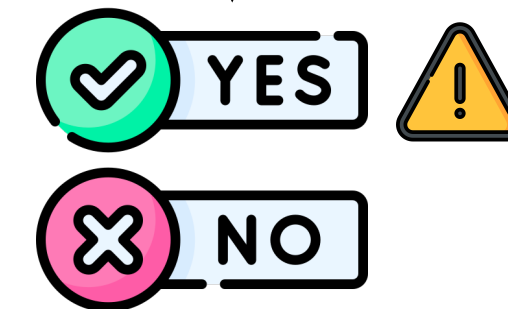
What is a filter?

Filter



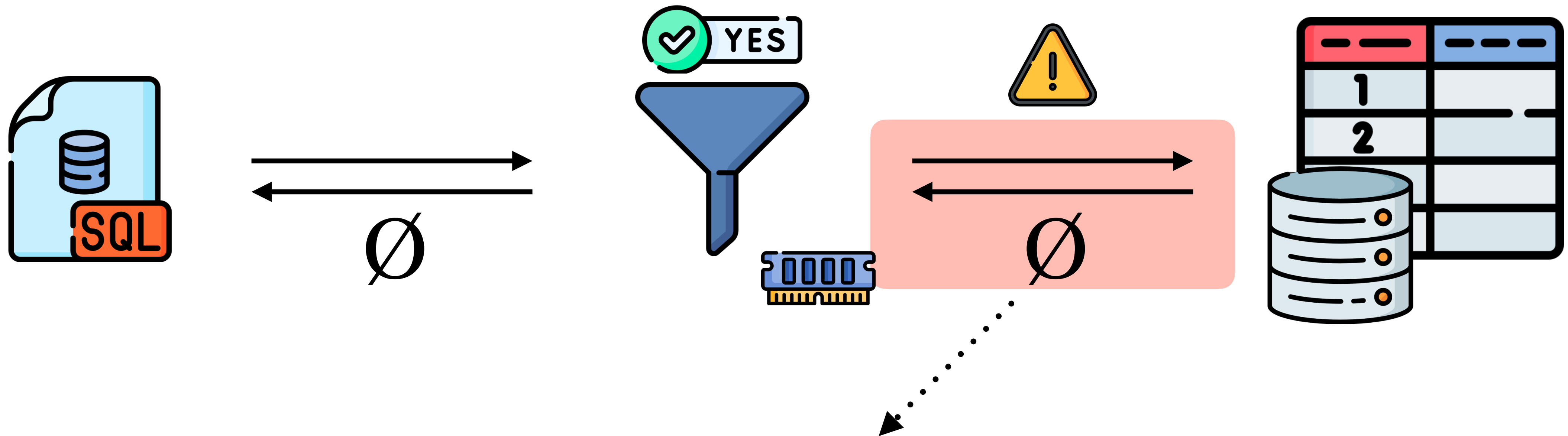
Answer membership queries with false positive probability ϵ

Does the dataset contain this key?



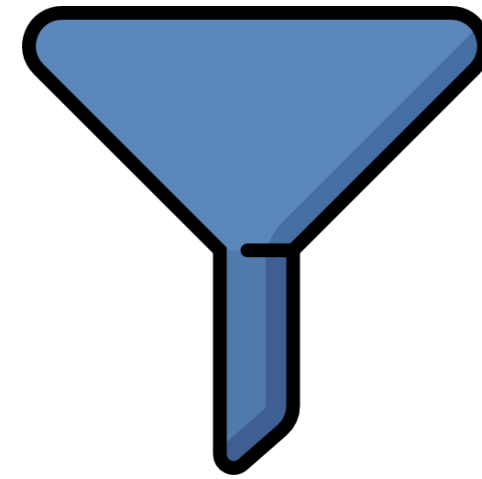
Non-existent key specified as
present with probability ϵ
 $0 < \epsilon < 1$

The cost of false positives



False positives incur disk I/O for non-existent keys

Unbounded false positive guarantees of filters



Is **X** in set? YES (**False positive**)

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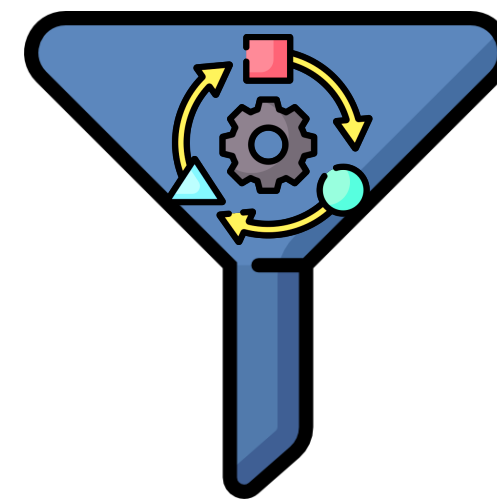
Time

Uniform Random: Close to $\epsilon \cdot n$ false positives w.h.p

Skewed workloads: Unbounded false positives, **up to n in worst case**

ϵ : false positive probability, n : number of queries

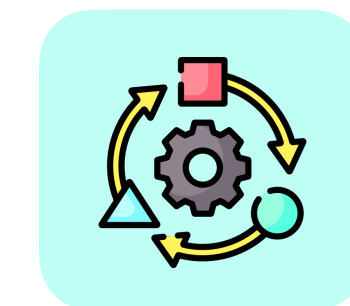
Bounding false positives with adaptivity^[BFG+'18]



Is **X** in set? YES (**False positive**)

Is **X** in set? NO (**True negative**)

Is **X** in set? NO (**True negative**)

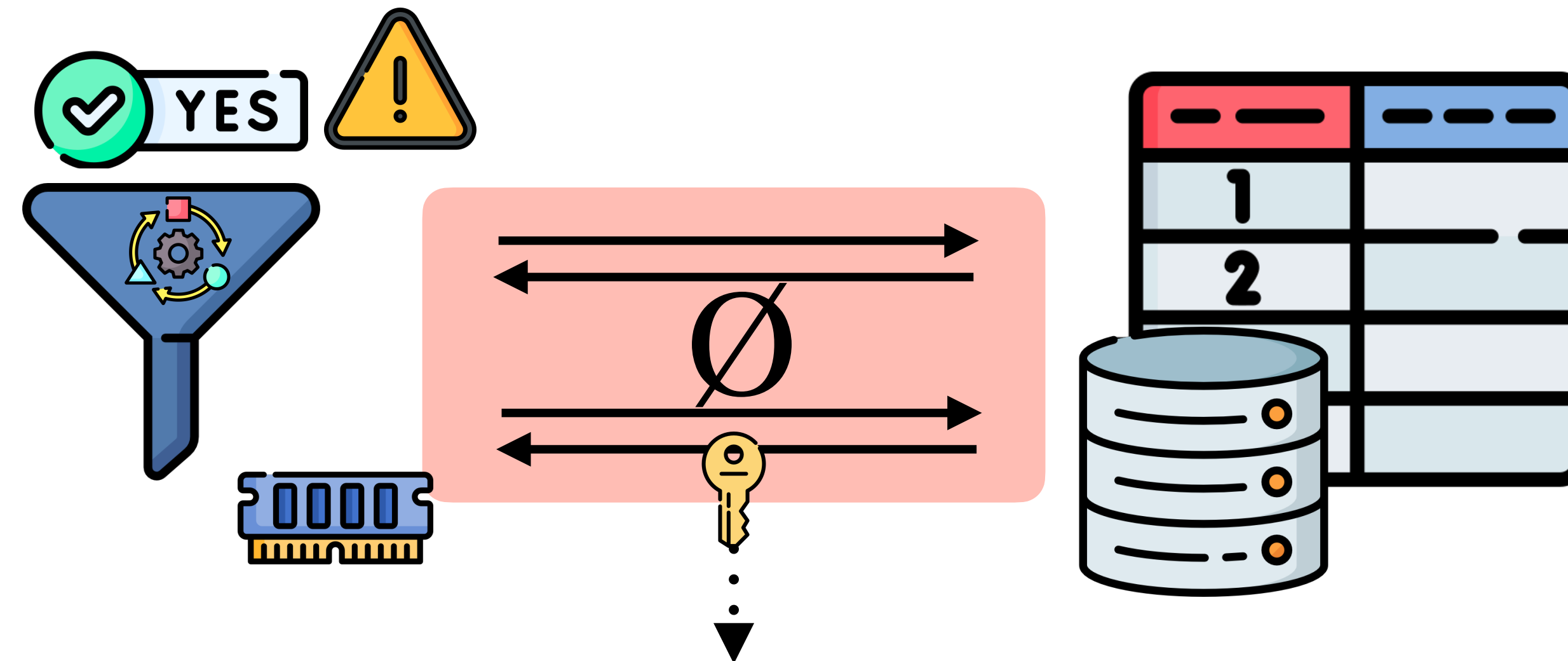


Time



Any Distribution: Close to $\epsilon \cdot n$ false positives w.h.p

Adaptive filters increase false positive cost

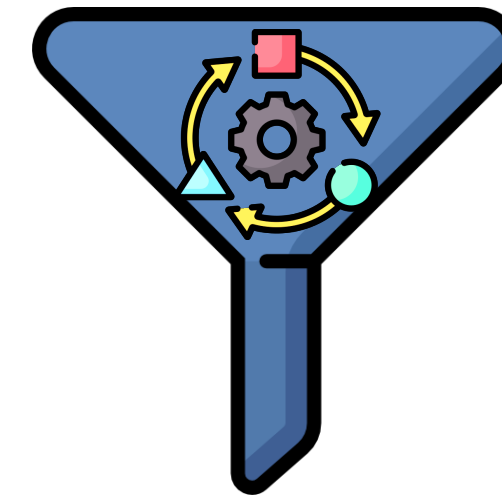
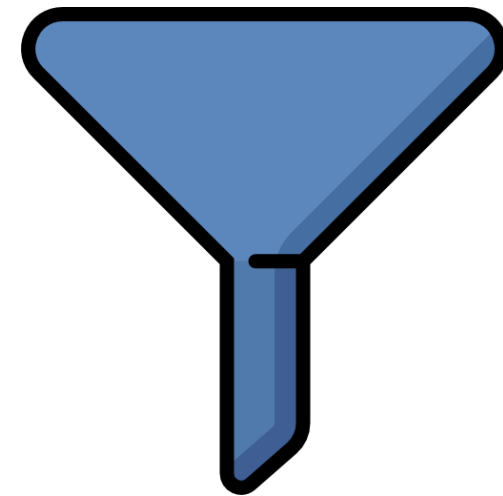


Adaptations require access to original keys, requiring additional disk accesses

Disk I/O Cost: $(1 + \alpha)$

α : relative extra I/O to fix false positive

False positive disk I/Os across distributions



Disk I/O cost

1

$(1 + \alpha)$

Uniform random

ϵn

$(1 + \alpha) \cdot (\epsilon n)$

Skewed

$\leq n$

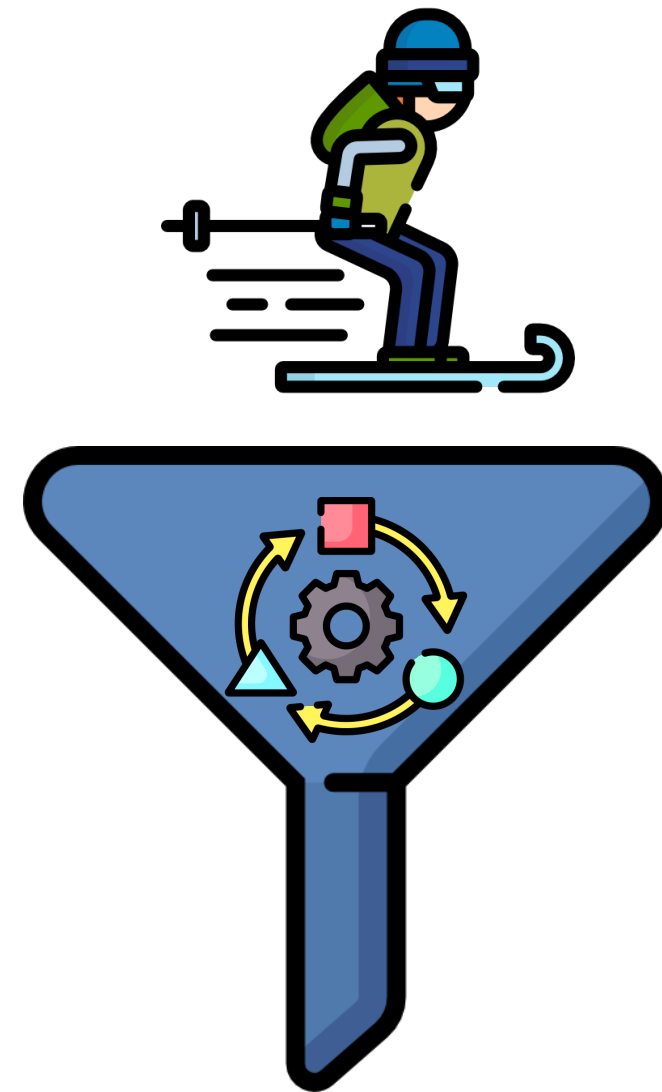
$(1 + \alpha) \cdot (\epsilon n)$

α : relative extra I/O to fix false positive

ϵ : false positive probability of single query

n : number of queries

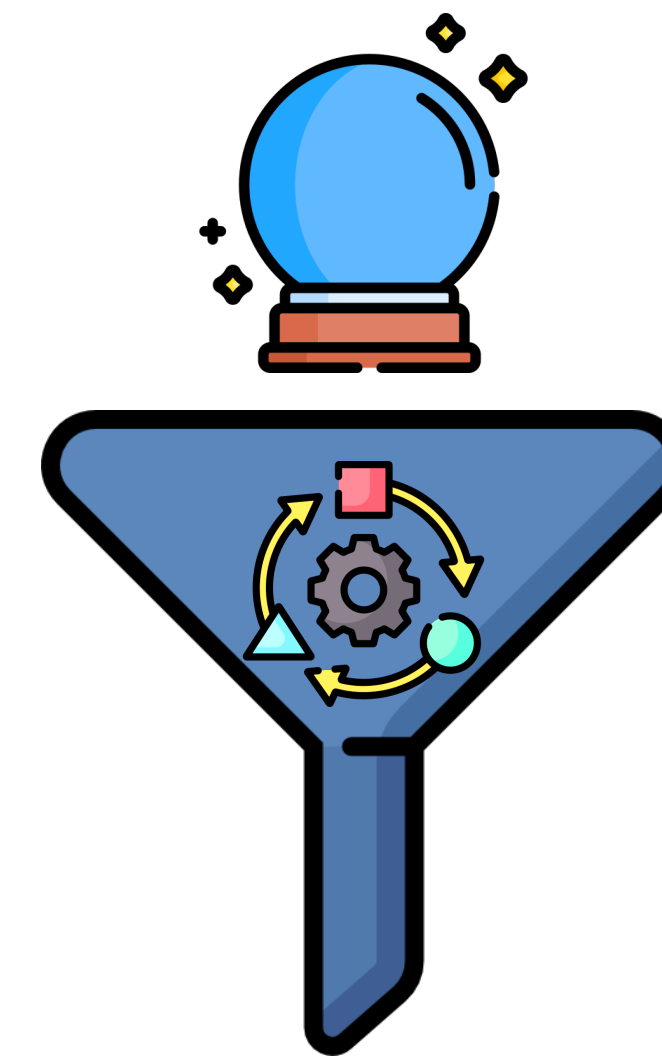
Online adaptivity: SkiQF and HybridSkiQF



SkiQF

Adaptivity as
ski rental problem

1.21x higher throughput over
AdaptiveQF across all workloads



HybridSkiQF

Periodically predicts if
adaptivity is required

1.28x higher throughput over
AdaptiveQF on stable workloads

The ski-rental problem

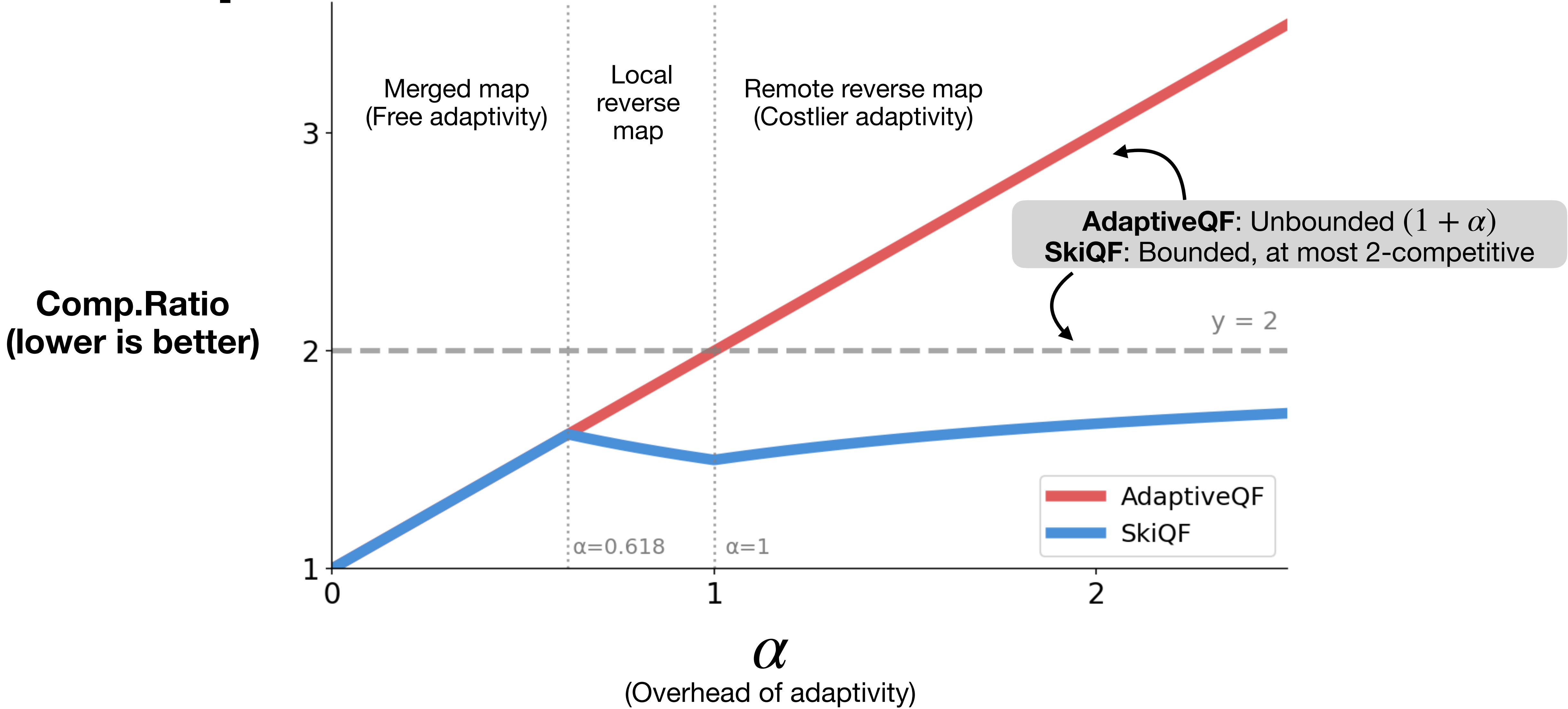


Buy or rent skis?

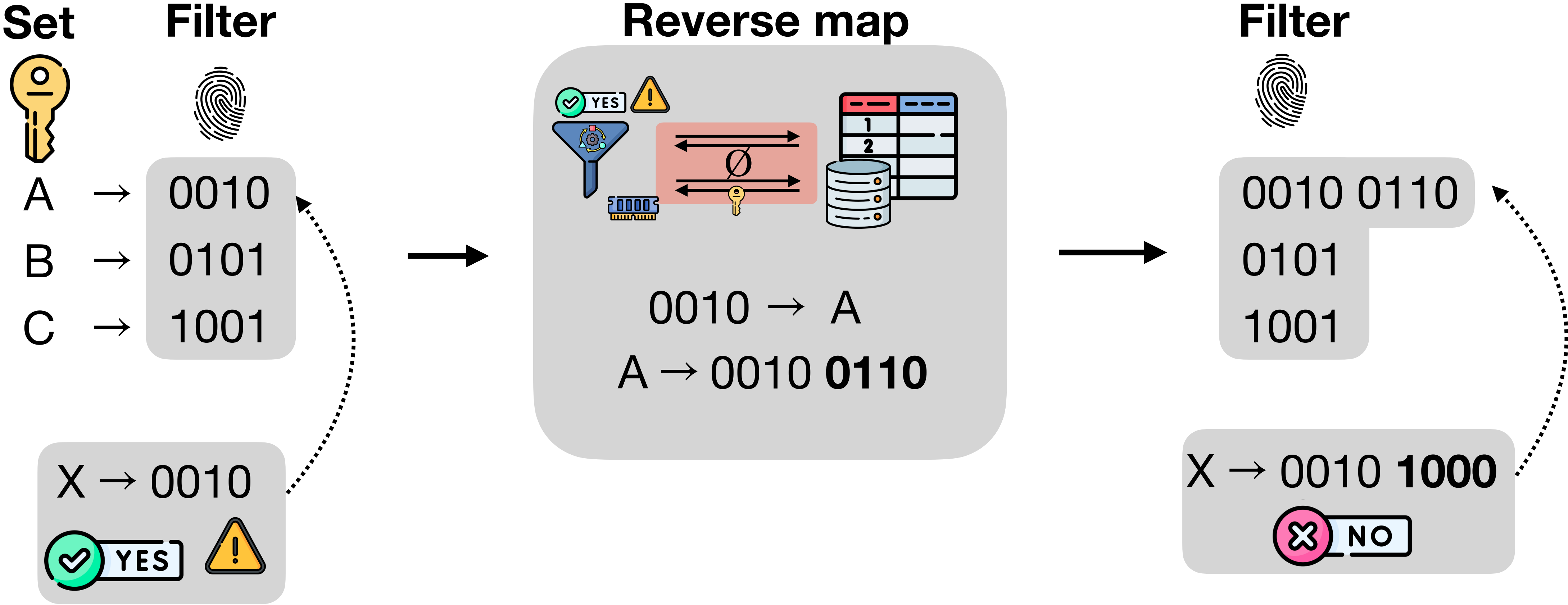
Algorithm : Keep renting until cumulative rental cost exceeds buy cost

Guarantees **2-competitiveness**: Never pays more than 2x the optimal offline cost

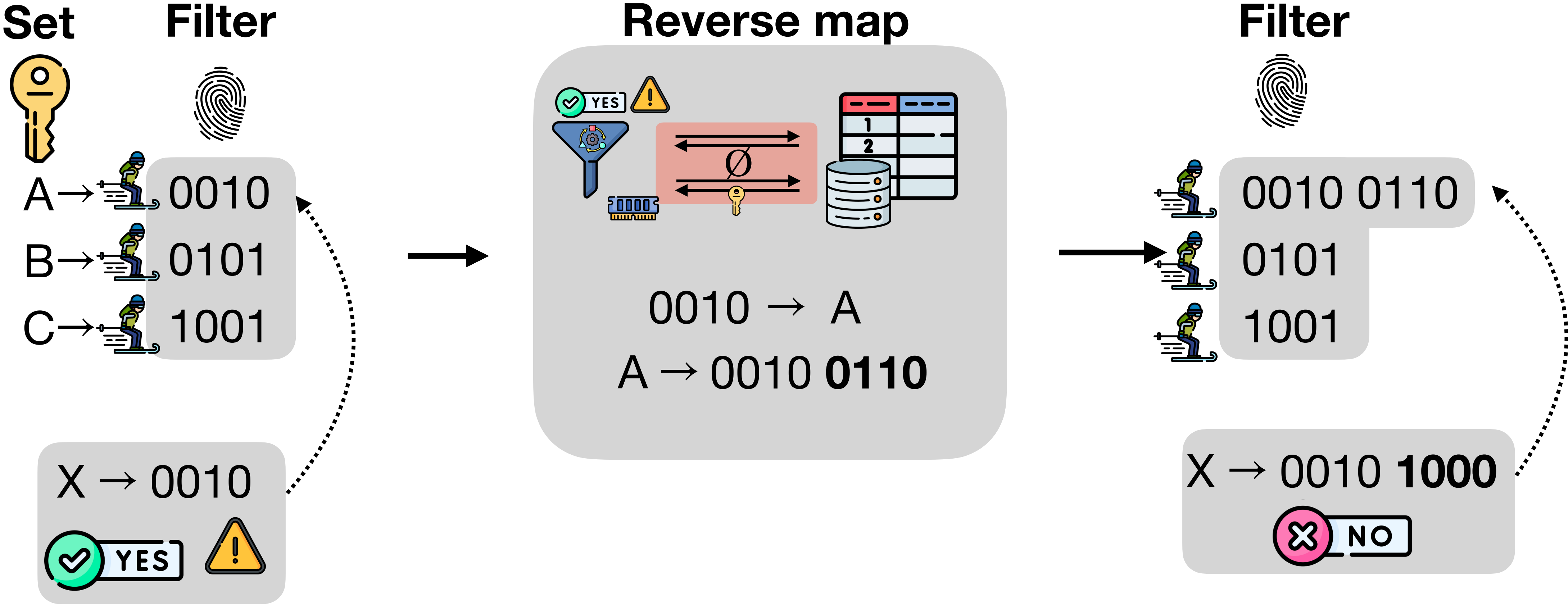
Competitiveness of SkiQF



How adaptations work



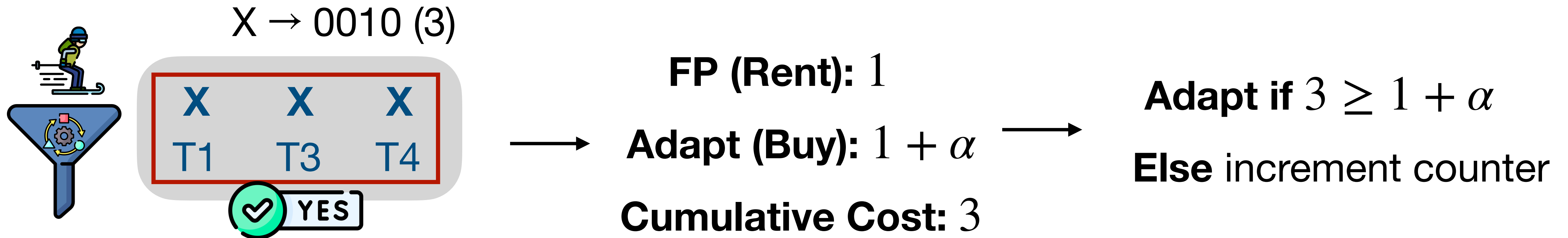
SkiQF: One ski-rental problem per fingerprint



Implementing the ski-rental algorithm

Algorithm : Keep renting until cumulative rental cost exceeds buy cost

Only requires quantifying buy, rent costs and counting prior occurrences



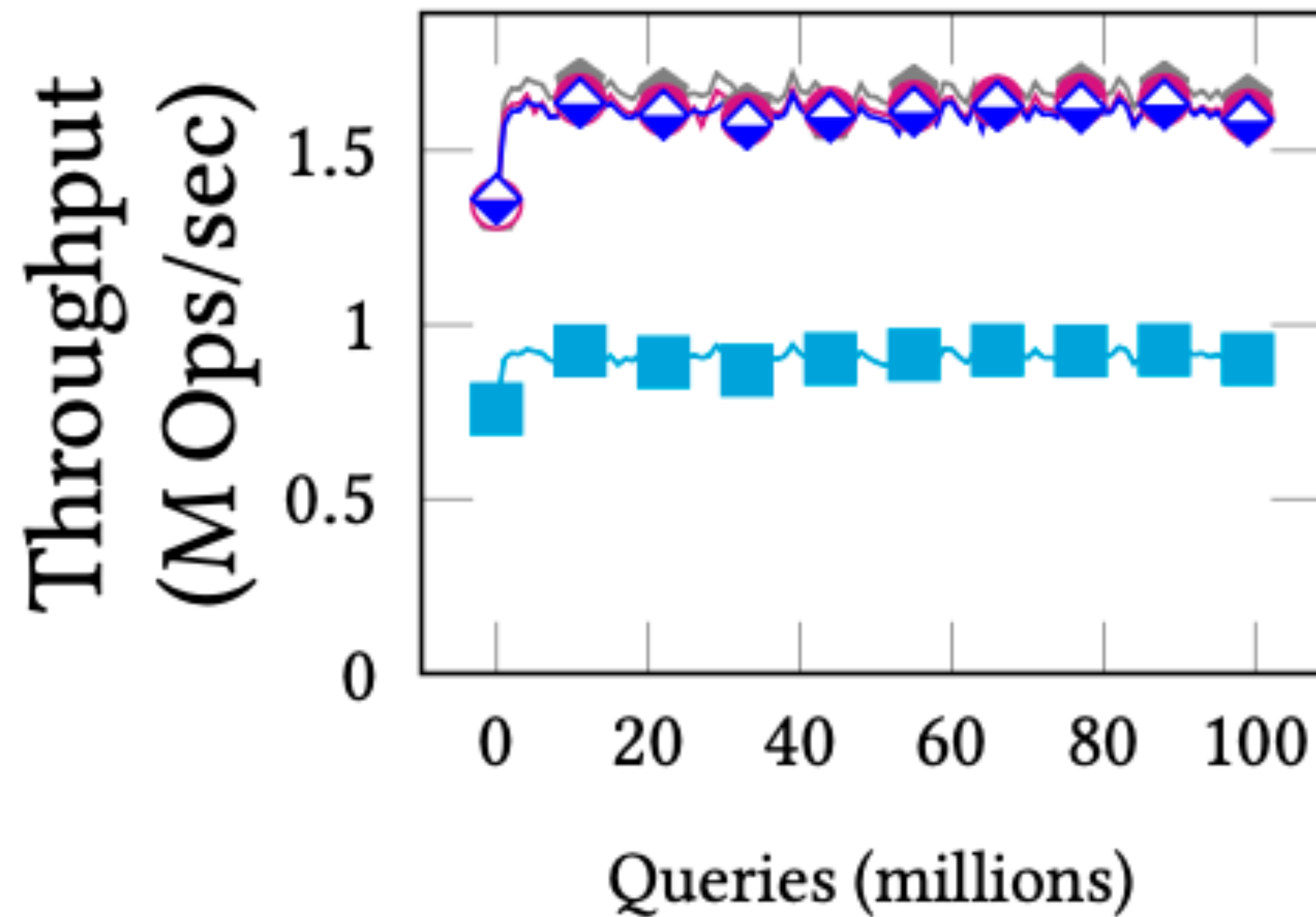
Experimental results

- WiredTiger database (~134 million keys, 100 million queries)
- $\alpha = 1$ (Adapting costs twice, paper also tests varying α)
- Baselines:
 - **Non-Adaptive:** Counting quotient filter (CQF)
 - **Adaptive:** AdaptiveQF
 - **Online Adaptive:** SkiQF, HybridSkiQF

Uniform random workload



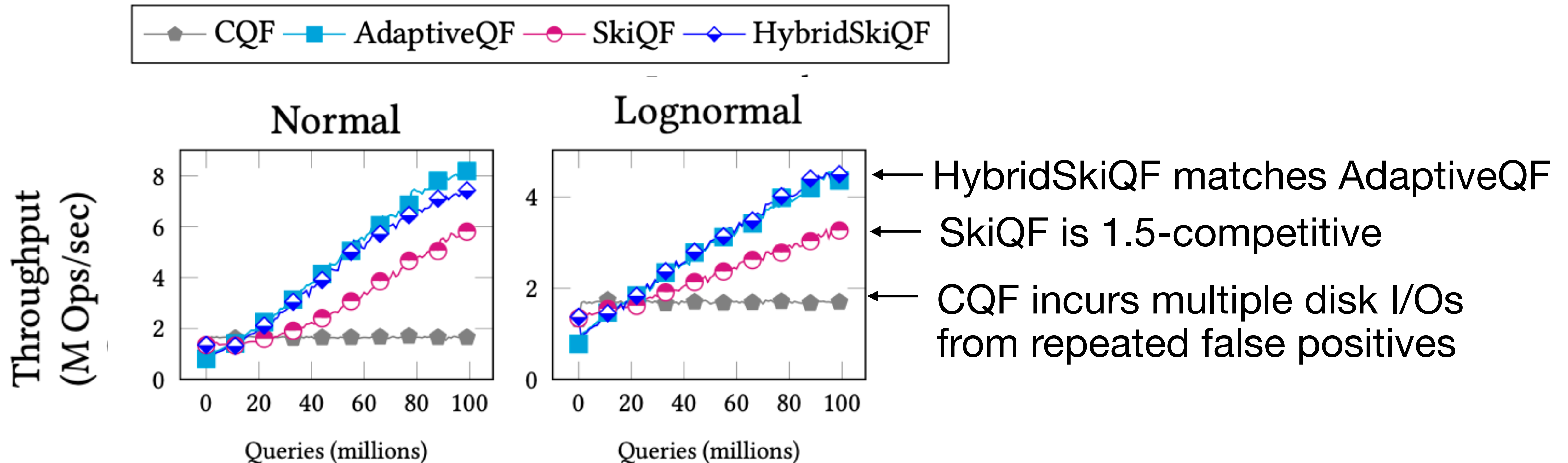
Uniform



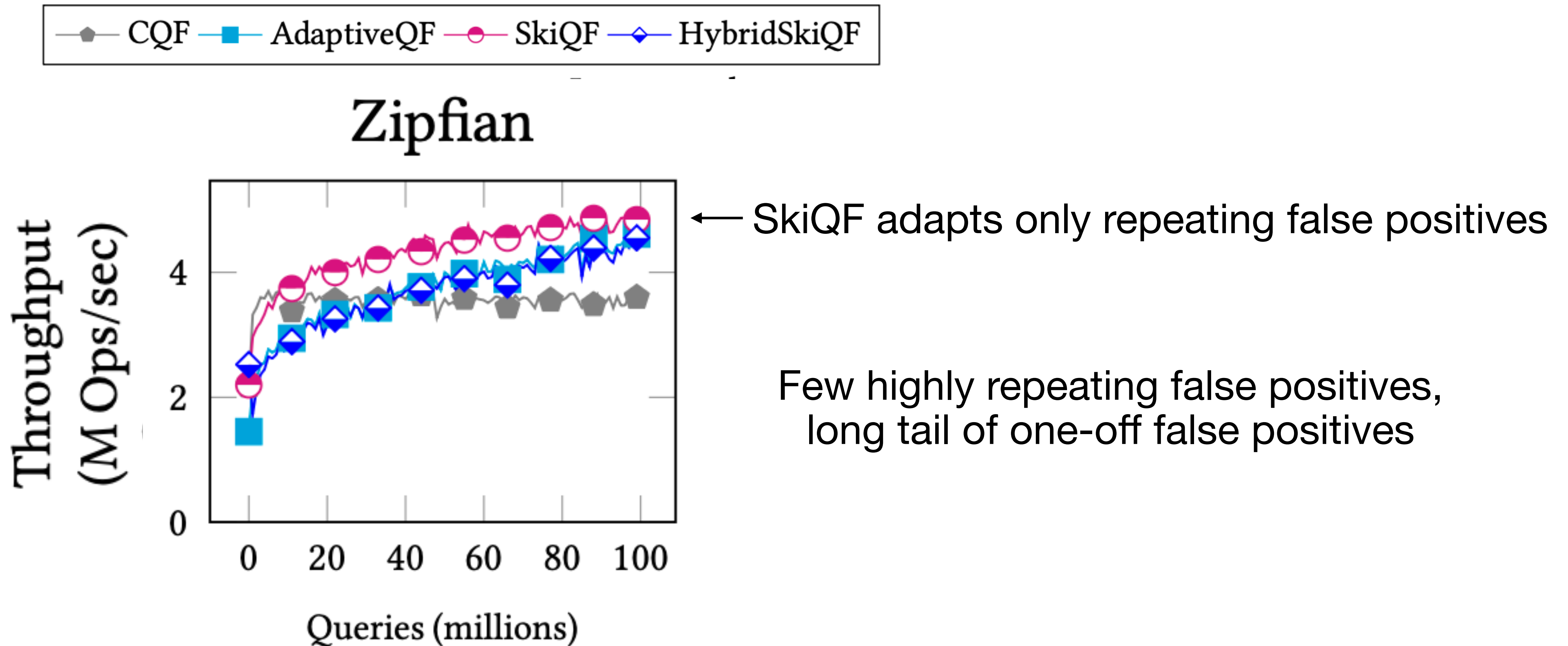
← Online adaptive filters match CQF

← AdaptiveQF has half the throughput of online-adaptive and non-adaptive filters

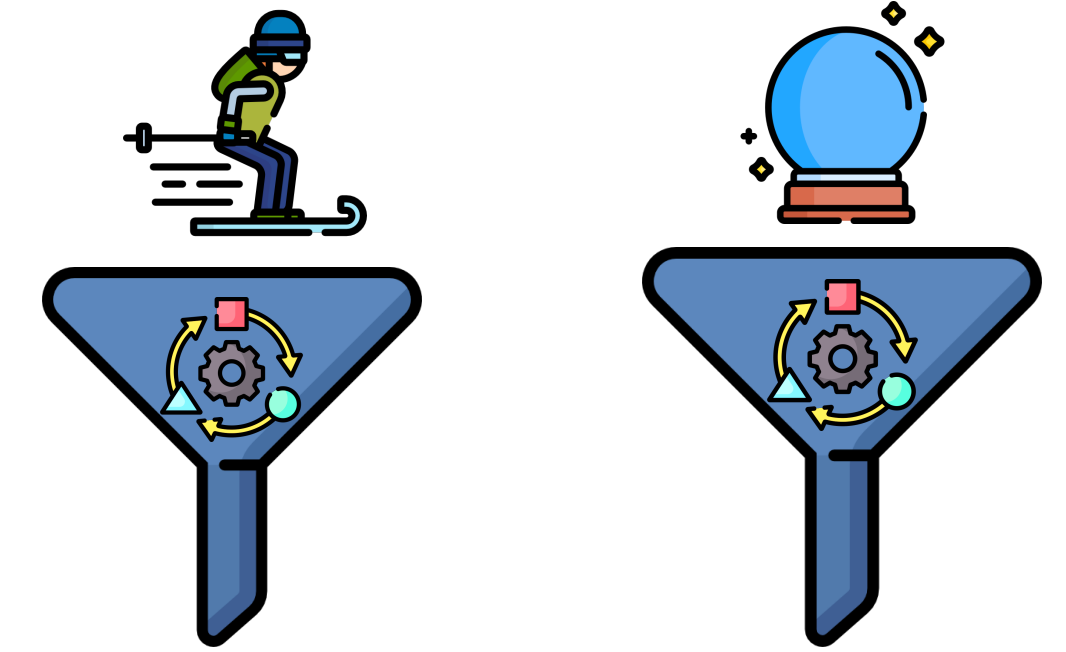
Skewed workloads



Zipfian workload (extreme skew)



Key takeaways



- All filters should be online adaptive!
- Online adaptivity:
 - **SkiQF**: Robust, bounded competitive performance for any workload
 - **HybridSkiQF**: Detects workload pattern, matches optimal filter for stable workloads
- **Easy to deploy**: Zero CPU/Memory overhead over AdaptiveQF
- **Code**: <https://github.com/saltsystemslab/SkiRentalAdaptiveQF>

